

Roles of energy storage in balancing power networks

(Joint work with Richard Gibbens and James Cruise)

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Outline

Introduction

Storage for arbitrage

Stochastic extensions

Integration of buffering for uncertainty

Demand-side management

Problem

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Hence there are major problems both of **variability** and of **uncertainty**.

Uses of storage

- ▶ Smoothing of **variability** — natural to address through **arbitrage**, **buy** energy when price is **low**, and **sell** energy when price is **high**.
- ▶ Buffering and control of **uncertainty** — can **integrate** with above by valuing energy in store.

Many other uses of **storage**, but we shall concentrate on the above two.





Dinorwig: capacity: 9 GWh rate: 1.8 GW efficiency 0.75–0.80

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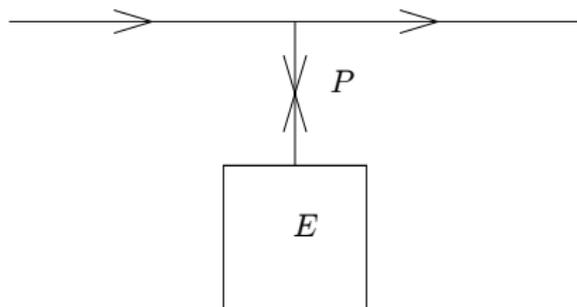
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Storage for arbitrage

Assume (initially)

- ▶ The **value** of storage can be captured in price **arbitrage** (buy energy when cheap and sell when expensive)
- ▶ Prices are known in advance (no uncertainty)

Model

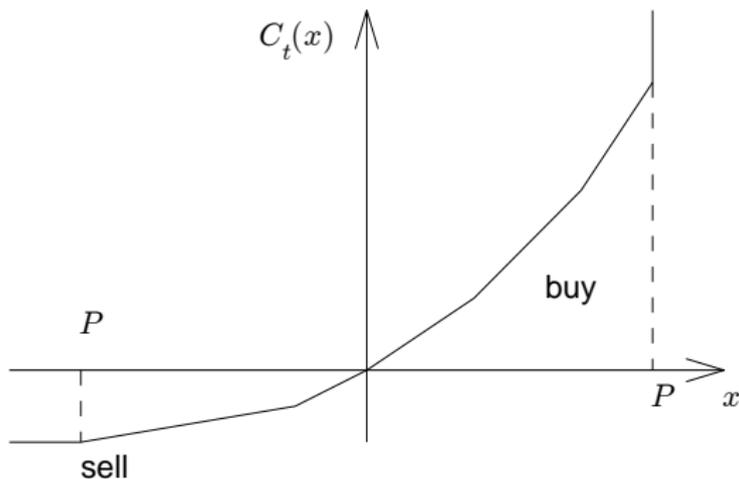


E = size of store — capacity constraint

P = max input/output rate — rate constraint

Cost function

At any (discrete) time t ,



$C_t(x)$ = cost of increasing level of store by x (positive or negative)
Assume **convex** (reasonable). This may model

- ▶ market impact
- ▶ efficiency of store
- ▶ rate constraints

Problem

Let $S_t =$ level of store at time t , $0 \leq t \leq T$.

Policy $S = (S_0, \dots, S_T)$, $S_0 = S_0^*$ (fixed), $S_T = S_T^*$ (fixed).

Define also $x_t(S) = S_t - S_{t-1}$

(energy “bought” by store at time t – positive or negative)

Problem: minimise cost

$$\sum_{t=1}^T C_t(x_t(S))$$

subject to

$$S_0 = S_0^*, \quad S_T = S_T^*$$

and

$$0 \leq S_t \leq E, \quad 1 \leq t \leq T - 1.$$

Result: Lagrangian sufficiency

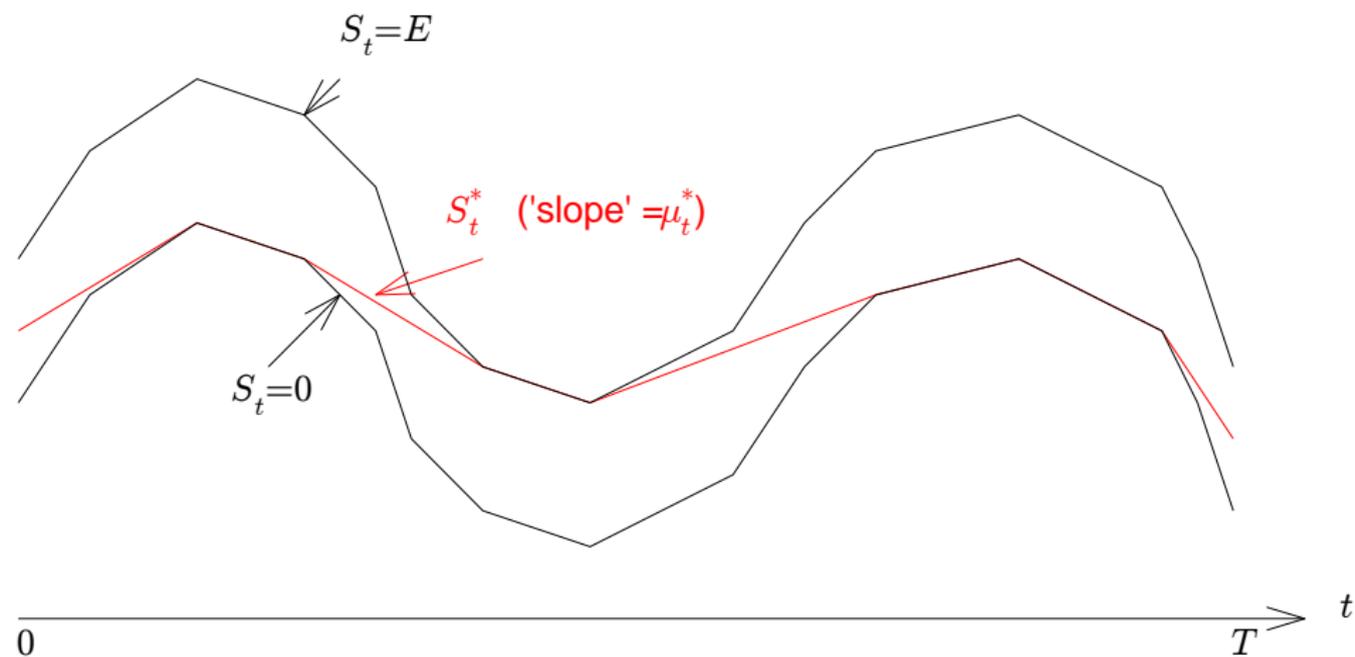
Suppose there exists a **vector** $\mu^* = (\mu_1^*, \dots, \mu_T^*)$ and a **value** $S^* = (S_0^*, \dots, S_T^*)$ of S such that

- ▶ S^* is **feasible**,
- ▶ $x_t(S^*)$ **minimises** $C_t(x) - \mu_t^* x$ over all x , $1 \leq t \leq T$,
- ▶ the pair (S^*, μ^*) satisfies the **complementary slackness conditions**, for $1 \leq t \leq T - 1$,

$$\begin{cases} \mu_{t+1}^* = \mu_t^* & \text{if } 0 < S_t^* < E, \\ \mu_{t+1}^* \leq \mu_t^* & \text{if } S_t^* = 0, \\ \mu_{t+1}^* \geq \mu_t^* & \text{if } S_t^* = E. \end{cases} \quad (1)$$

Then S^* **solves** the stated problem.

Picture proof



Proper proof

Let S be any vector which is **feasible** for the problem (with $S_0 = S_0^*$ and $S_T = S_T^*$). Then, from the **minimisation condition** above,

$$\sum_{t=1}^T [C_t(x_t(S^*)) - \mu_t^* x_t(S^*)] \leq \sum_{t=1}^T [C_t(x_t(S)) - \mu_t^* x_t(S)].$$

Rearranging and recalling that S and S^* agree at 0 and at T , we have

$$\begin{aligned} \sum_{t=1}^T C_t(x_t(S^*)) - \sum_{t=1}^T C_t(x_t(S)) &\leq \sum_{t=1}^T \mu_t^* (S_t^* - S_{t-1}^* - S_t + S_{t-1}) \\ &= \sum_{t=1}^{T-1} (S_t^* - S_t) (\mu_t^* - \mu_{t+1}^*) \\ &\leq 0, \end{aligned}$$

by the **complementary slackness conditions** above.

Comment

- ▶ The above result (essentially an application of the Lagrangian Sufficiency Principle) does **not** require **convexity** of the functions C_t .
- ▶ However, **convexity** of the functions C_t is sufficient to **guarantee** the **existence** of a pair (S^*, μ^*) as above.

The latter result is a standard application of **Strong Lagrangian** theory (i.e. the **Supporting Hyperplane Theorem**).

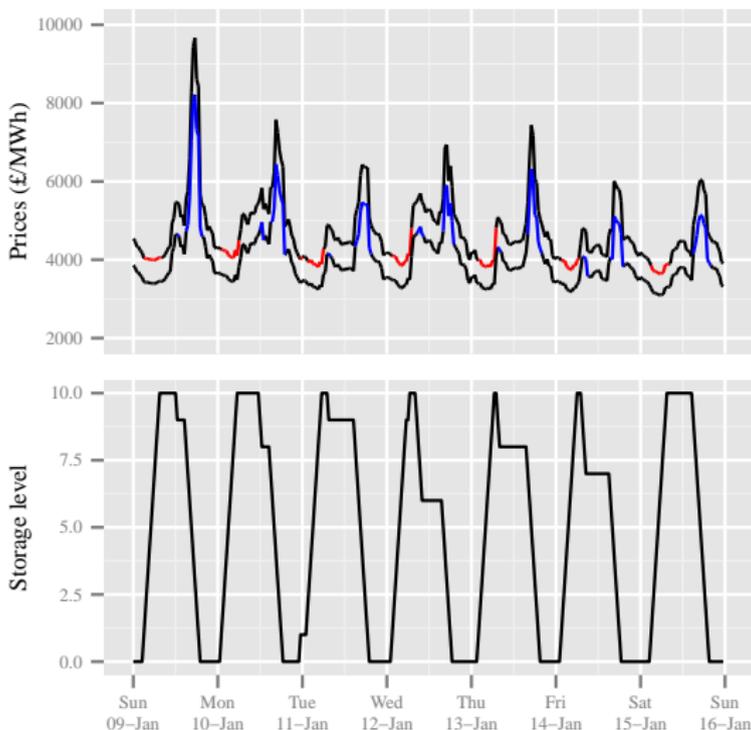
Algorithm

- ▶ Need to **identify** the relevant value μ_t^* at each time t .
- ▶ This is as in the previous picture (“**stretched string**”).
- ▶ Note that, at each time t , only a **typically short time horizon** is required to identify μ_t^* .

Example: Real prices with Dinorwig parameters

$E/P = 5$ hrs Efficiency = 0.85 (ratio of sell to buy price).

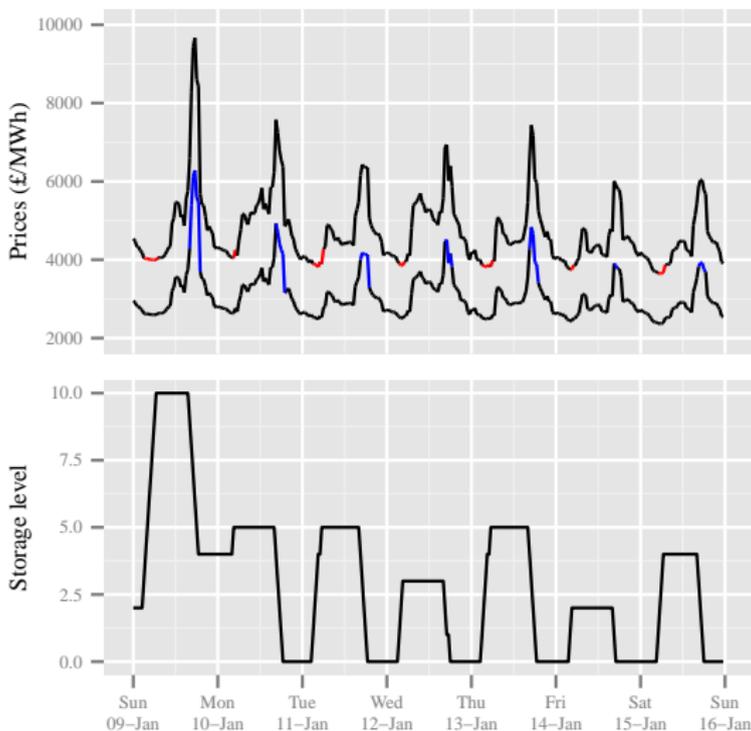
Solution is bang-bang: red points buy, blue points sell



Example: Real prices with Dinorwig parameters

$E/P = 5$ hrs Efficiency = 0.65 (ratio of sell to buy price).

Solution is bang-bang: red points buy, blue points sell



Small store

This is a store whose activities are not so great as to impact upon the market, and which thus has linear buy and sell prices. Thus, for all t ,

$$C_t(x) = \begin{cases} c_t^{(b)} x & \text{if } 0 \leq x \leq P \\ c_t^{(s)} x & \text{if } -P \leq x < 0 \end{cases}$$

where $0 < c_t^{(s)} \leq c_t^{(b)}$ and P is rate constraint.

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Characteristics

- ▶ Optimal control is bang-bang: at each time buy as much as possible, do nothing, or sell as much as possible (according to value of μ_t^* in relation to above costs).
- ▶ If $E = \infty$ (no capacity constraint) then optimal solution is global in time.
- ▶ If $P = \infty$ (no rate constraint) then optimal solution is very local in time.

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Result

Suppose that the cost functions C_t evolve **randomly** in time, and that

$$C_t = \xi_t \bar{C}_t, \quad 1 \leq t \leq T,$$

where $(\bar{C}_1, \dots, \bar{C}_T)$ is a sequence of **deterministic** cost functions and where (ξ_1, \dots, ξ_T) is a sequence of strictly positive real-valued random variables such that

$$\mathbf{E}(\xi_t | \mathcal{F}_{t-1}) = \xi_{t-1}, \quad 1 \leq t \leq T,$$

and where the deterministic functions \bar{C}_t are assumed to satisfy the same conditions as the cost functions C_t of the deterministic problem.

Then the **optimal control** (cost minimisation) strategy remains **exactly as previously**, with the optimal sequence of store levels as given in the case where **stochastic cost functions** C_t are replaced by their **deterministic counterparts** \bar{C}_t .

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Integration of buffering for uncertainty

Suppose that the store is **also used** to provide **buffering** against **relatively rare** unexpected events.

Then the previous **optimization problem** may be rewritten as:

Problem: minimise

$$\sum_{t=1}^T [C_t(x_t(S)) + A_t(S_t)]$$

subject to

$$S_0 = S_0^*, \quad S_T = S_T^*$$

and

$$0 \leq S_t \leq E, \quad 1 \leq t \leq T - 1.$$

Here the functions A_t , such that $A_t(S_t)$ is the cost (negative benefit) of having a level of **reserve** S_t in the store at time t , are **independently calculable**.

Result - Lagrangian sufficiency

Assume, for simplicity, we incorporate the **capacity constraints** into the functions A_t , and the functions C_t and A_t are **differentiable**.

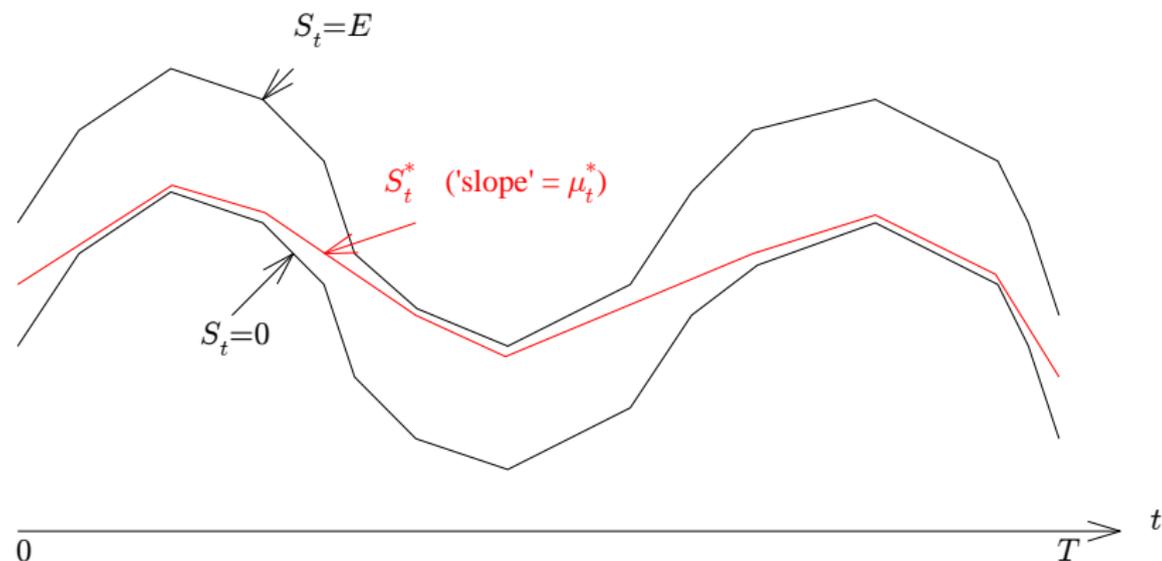
Suppose there exists a **scalar** ν^* and a **value** $S^* = (S_0^*, \dots, S_T^*)$ (with S_0^* and S_T^* having their required values) such that

$$C'_t(x_t) = \nu - \sum_{u=t}^T A'_u(S_u) \quad =: \mu_t$$

Then S^* **solves** the stated problem.

Lagrangian sufficiency (ctd)

The solution has a “stretched string” characterisation as before:



The previous algorithm may be adapted to give a solution which is essentially **local in time**.

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Like **storage**, this is **shifting energy through time**.

However, there are typically **additional constraints** e.g.

- ▶ **duration of actions**
- ▶ **frequency of actions.**